

Discussion

T-49-19-07

Super8™ MCU ROMIess, ROM, and Prototyping Device with EPROM Interface

Z8800, Z8801, Z8820, Z8822

FEATURES

- Improved Z8[®] instruction set includes multiply and divide instructions, Boolean and BCD operations.
- Additional instructions support threaded-code languages, such as "Forth."
- 325 byte registers, including 272 general-purpose registers, and 53 mode and control registers.
- Addressing of up to 128K bytes of memory.
- Two register pointers allow use of short and fast instructions to access register groups within 600 nsec.
- Direct Memory Access controller (DMA).
- Two 16-bit counter/timers.

- Up to 32 bit-programmable and 8 byte-programmable I/O lines, with 2 handshake channels.
- Interrupt structure supports:
 - ☐ 27 interrupt sources
 - ☐ 16 Interrupt vectors (2 reserved for future versions)
 - □ 8 interrupt levels
 - ☐ Servicing in 600 nsec. (1 level only)
- Full-duplex UART with special features.
- On-chip oscillator.
- 20 MHz clock.
- 8K byte ROM for Z8820

GENERAL DESCRIPTION

The Zilog Super8 single-chip MCU can be used for development and production. It can be used as I/O- or memory-intensive computers, or configured to address external memory while still supporting many I/O lines.

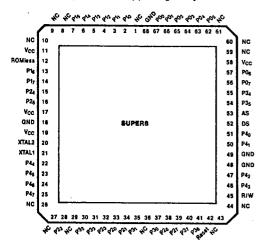
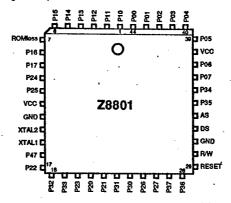


Figure 1a. Pin Assignments - 68-pin PLCC

The Super8 features a full-duplex universal asynchronous receiver/transmitter (UART) with on-chip baud rate generator, two programmable counter/timers, a direct memory access (DMA) controller, and an on-chip oscillator.

The Super8 is also available as a 48-pin and 68-pin ROMless microcomputer with four byte-wide I/O ports plus a byte-wide address/data bus. Additional address bits can be configured, up to a total of 16.



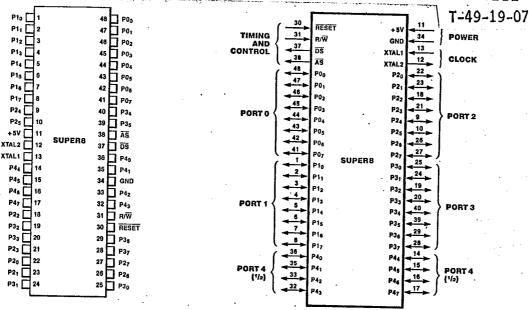


Figure 1b. Pin Assignments - 48-pin DIP

Figure 2. Pin Functions

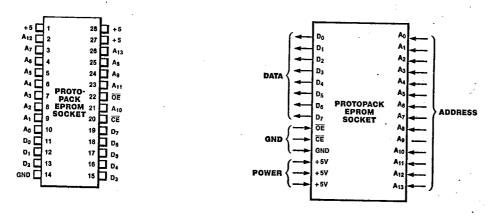


Figure 3. Pin Assignments—28-Pin Piggyback Socket

Figure 4. Pin Functions—28-Pin Piggyback Socket

Protopack

This part functions as an emulator for the basic microcomputer. It uses the same package and pin-out as the basic microcomputer but also has a 28-pin "piggy back" socket on the top into which a ROM or EPROM can be installed. The socket is designed to accept a type 2764 EPROM.

This package permits the protopack to be used in prototype and final PC boards while still permitting user program

development. When a final program is developed, it can be mask-programmed into the production microcomputer device, directly replacing the emulator. The protopack part is also useful in situations where the cost of mask-programming is prohibitive or where program flexibility is desired.

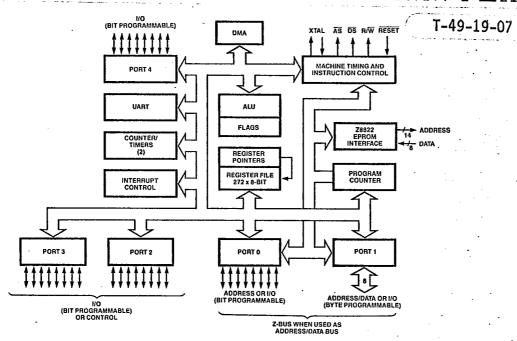


Figure 5. Functional Block Diagram

ARCHITECTURE

The Super8 architecture includes 325 byte-wide internal registers. 272 of these are available for general purpose use; the remaining 53 provide control and mode functions.

The instruction set is specially designed to deal with this large register set. It includes a full complement of 8-bit arithmetic and logical operations, including multiply and divide instructions and provisions for BCD operations. Addresses and counters can be incremented and decremented as 16-bit quantities. Rotate, shift, and bit manipulation instructions are provided. Three new instructions support threaded-code languages.

The UART is a full-function multipurpose asynchronous serial channel with many premium features.

The 16-bit counters can operate independently of be cascaded to perform 32-bit counting and timing operations. The DMA controller handles transfers to and from the register file or memory. DMA can use the UART or one of two ports with handshake capability.

The architecture appears in the block diagram (Figure 5).

PIN DESCRIPTIONS

The Super8 connects to external devices via the following TTL-compatible pins:

 $\overline{\text{AS.}}$ Address Strobe (output, active Low). $\overline{\text{AS}}$ is pulsed Low once at the beginning of each machine cycle. The rising edge indicates that addresses $\overline{\text{RW}}$ and $\overline{\text{DM}}$, when used, are valid.

DS. Data Strobe (output, active Low). \overline{DS} provides timing for data movement between the address/data bus and external memory. During write cycles, data output is valid at the leading edge of \overline{DS} . During read cycles, data input must be valid prior to the trailing edge of \overline{DS} .

P0₀-P0₇, P1₀-P1₇, P2₀-P2₇, P3₀-P3₇, P4₀-P4₇. Port I/O Lines (input/output). These 40 lines are divided into five 8-bit I/O ports that can be configured under program control for I/O or external memory interface.

In the ROMless devices, Port 1 is dedicated as a multiplexed address/data port, and Port 0 pins can be assigned as additional address lines; Port 0 non-address pins may be assigned as I/O. In the ROM and protopack, Port 1 can be assigned as input or output, and Port 0 can be assigned as input or output on a bit by bit basis.

Ports 2 and 3 can be assigned on a bit-for-bit basis as general I/O or interrupt lines. They can also be used as special-purpose I/O lines to support the UART, counter/timers, or handshake channels.

Port 4 is used for general I/O.

During reset, all port pins are configured as inputs (high impedance) except for Port 1 and Port 0 in the ROMless devices. In these, Port 1 is configured as a multiplexed address/data bus, and Port 0 pins P0₀-P0₄ are configured as address out, while pins P0₅-P0₇ are configured as inputs.

RESET. Reset (input, active Low). Reset initializes and starts the Super8. When it is activated, it halts all processing; when

it is deactivated, the Super8 begins processing at address 0020_H. T-49-19-07

ROMIess. (input, active High). This input controls the operation mode of a 68-pin Super8. When connected to V_{CC} , the part will function as a ROMIess Z8800. When connected to GND, the part will function as a Z8820 ROM part.

R/W. Read/Write (output). R/W determines the direction of data transfer for external memory transactions. It is Low when writing to program memory or data memory, and High for everything else.

XTAL1, XTAL2. (Crystal oscillator input.) These pins connect a parallel resonant crystal or an external clock source to the on-board clock oscillator and buffer.

REGISTERS

The Super8 contains a 256-byte internal register space. However, by using the upper 64 bytes of the register space more than once, a total of 325 registers are available.

Registers from 00 to BF are used only once. They can be accessed by any register command. Register addresses C0 to FF contain two separate sets of 64 registers. One set, called control registers, can only be accessed by register direct commands. The other set can only be addressed by register indirect, indexed, stack, and DMA commands.

The uppermost 32 register direct registers (E0 to FF) are further divided into two banks (0 and 1), selected by the Bank Select bit in the Flag register. When a Register Direct command accesses a register between E0 and FF, it looks at the Bank Select bit in the Flag register to select one of the banks.

The register space is shown in Figure 6.

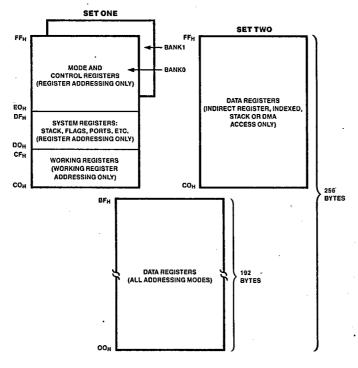


Figure 6. Super8 Registers

Working Register Window

Control registers R214 and R215 are the register pointers, RP0 and RP1. They each define a moveable, 8-register section of the register space. The registers within these spaces are called working registers.

Working registers can be accessed using short 4-bit addresses. The process, shown in section a of Figure 4, works as follows:

- The high-order bit of the 4-bit address selects one of the two register pointers (0 selects RP0; 1 selects RP1).
- The five high-order bits in the register pointer select an 8-register (contiguous) slice of the register space.
- The three low-order bits of the 4-bit address select one of the eight registers in the slice.

a. 4-Bit Addressing

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The net effect is to concatenate the five bits from the register pointer to the three bits from the address to form an 8-bit address. As long as the address in the register pointer remains unchanged, the three bits from the address will always point to an address within the same eight registers.

The register pointers can be moved by changing the five high bits in control registers R214 for RP0 and R215 for RP1.

The working registers can also be accessed by using full 8-bit addressing. When an 8-bit logical address in the range 192 to 207 (C0 to CF) is specified, the lower nibble is used similarly to the 4-bit addressing described above. This is shown in section b of Figure 7.

b. 8-Bit Addressing

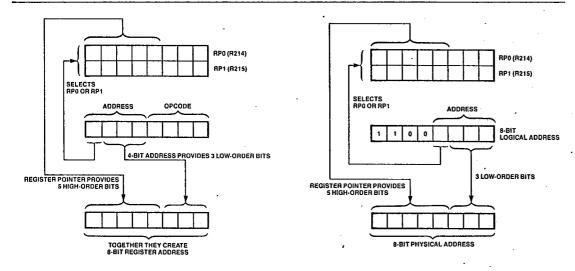


Figure 7. Working Register Window

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Since any direct access to logical addresses 192 to 207 involves the register pointers, the physical registers 192 to 207 can be accessed only when selected by a register pointer. After a reset, RP0 points to R192 and RP1 points to R200.

Register List T-49-19-07

Table 1 lists the Super8 registers. For more details, see Figure 8.

Table 1. Super-8 Registers

Addre			
Decimal	Hexadecimal	Mnemonic	Function
eneral-Purpose Reg	isters	,	
000-192	00-BF		General purpose (all address modes)
192-207	C0-CF	- .	Working register (direct only)
192-255	C0-FF		General purpose (indirect only)
ode and Control Re	gisters		
208	D0	P0	Port 0 I/O bits
209	D1	, P1	Port 1 (I/O only)
210	D2	P2	Port 2
211	D3	P3	Port 3
212	D4	P4	Port 4
213	D5 .	FLAGS	System Flags Register
214	D6	RP0	Register Pointer 0
215	D7	RP1	Register Pointer 1
216	D8	SPH	Stack Pointer High Byte
217	D9	SPL	Stack Pointer Low Byte
218	DA	IPH	Instruction Pointer High Byte
219	DB	IPL	Instruction Pointer Low Byte
220	DC	IRQ	Interrupt Request
221	DĎ	IMR	• • •
222	DE	SYM	Interrupt Mask Register
222 224			System Mode
224	E0 Bank 0	COCT	. CTR 0 Control
005	Bank 1	COM	CTR 0 Mode
225	E1 Bank 0	C1CT	CTR 1 Control
	Bank 1	C1M .	CTR 1 Mode
226	E2 Bank 0	, COCH	CTR 0 Capture Register, bits 8-15
	Bank 1	CTCH	CTR 0 Timer Constant, bits 8-15
227	E3 Bank 0	COCL ,	CTR 0 Capture Register, bits 0-7
š.	Bank 1	CTCL	CTR 0 Time Constant, bits 0-7
228	E4' Bank 0	C1CH	CTR 1 Capture Register, bits 8-15
	Bank 1	C1TCH	CTR 1 Time Constant, bits 8-15
229	E5 Bank 0	C1CL	CTR 1 Capture Register, bits 0-7
	Bank 1	C1TCL	CTR 1 Time Constant, bits 0-7
235	EB Bank 0	´UTC	UART Transmit Control
236	EC Bank 0	URC	UART Receive Control
237	ED Bank 0	UIE	UART Interrupt Enable
239	EF Bank 0	UIO	UART Data
240	F0 Bank 0	POM	Port 0 Mode
	Bank 1	DCH	DMA Count, bits 8-15
241	F1 Bank 0	PM	Port Mode Register
	Bank 1	DCL	DMA Count, bits 0-7
244	. F4 Bank 0	HOC	Handshake Channel 0 Control
245	F5 Bank 0	H1C	Handshake Channel 1 Control
246	F6 Bank 0	P4D	Port 4 Direction
247	F7 Bank 0	P4OD	
248			Port 4 Open Drain
440		. P2AM	Port 2/3 A Mode
	Bank 1	UBGH .	UART Baud Rate Generator, bits 8-15

Table 1. Super-8 Registers (Continued)

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	-4		-1		_	u	•

Addr	ess		
Decimal	Hexadecimal	Mnemonic	Function
lode and Control Re	gisters (Continued)		
249	F9 Bank 0	P2BM	Port 2/3 B Mode
	Bank 1	UBGL	UART Baud Rate Generator, bits 0-7
250	FA Bank 0	P2CM	Port 2/3 C Mode
	Bank 1	UMA ·	UART Mode A
251	FB Bank 0	P2DM	Port 2/3 D Mode
	Bank 1	UMB	UART Mode B
252	FC Bank 0	P2AIP	Port 2/3 A Interrupt Pending
253	FD Bank 0	P2BIP	Port 2/3 B Interrupt Pending
254	FE Bank 0	EMT	External Memory Timing
	Bank 1	WUMCH	Wakeup Match Register
255	. FF Bank 0	IPR	Interrupt Priority Register
	Bank 1	WUMSK	Wakeup Mask Register

MODE AND CONTROL REGISTERS

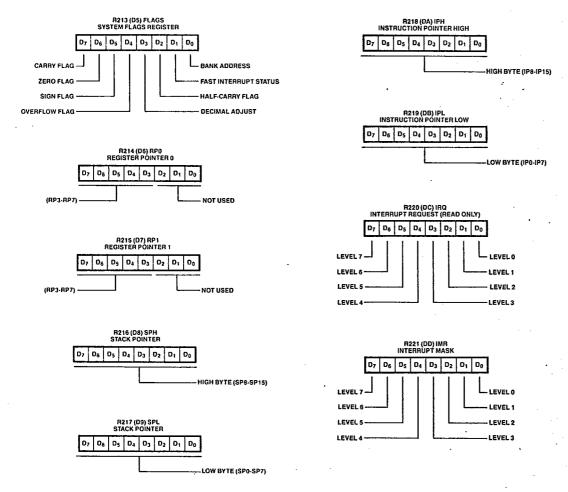
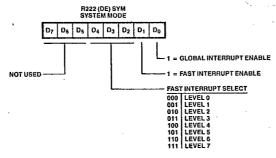
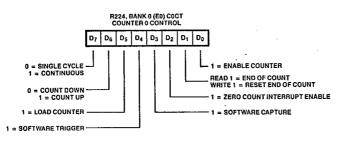


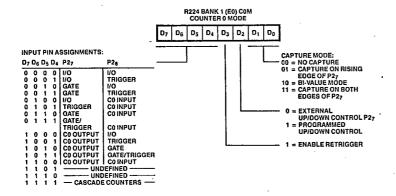
Figure 8. Mode and Control Registers

MODE AND CONTROL REGISTERS (Continued)

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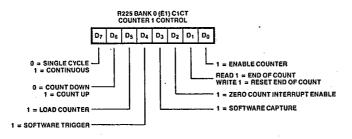
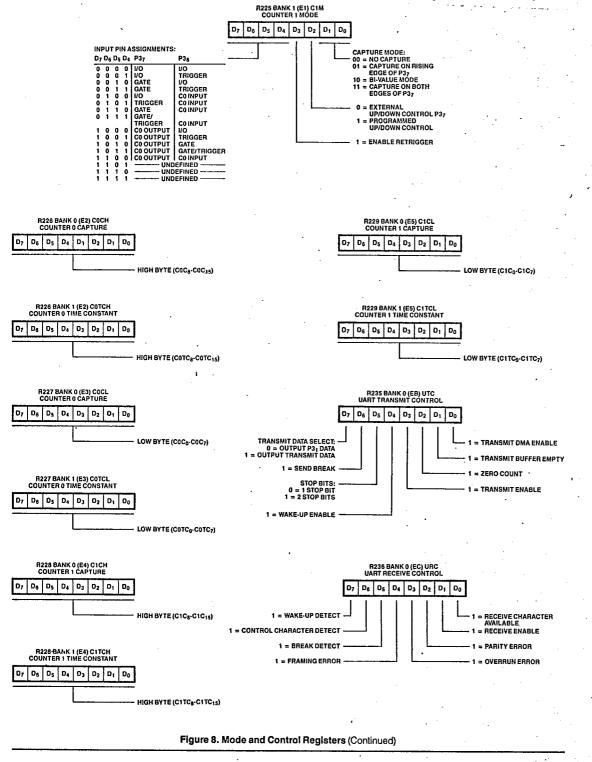


Figure 8. Mode and Control Registers (Continued)

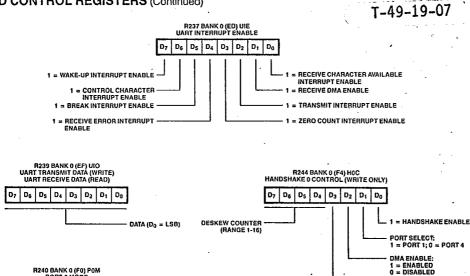
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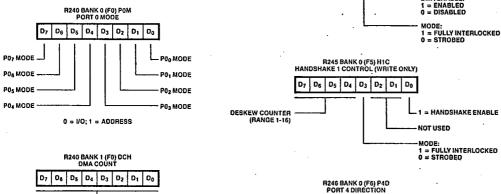


D6 D5 D4 D3 D2 D1 D0

P40-P47 I/O DIRECTION 0 = OUTPUT; 1 = INPUT

MODE AND CONTROL REGISTERS (Continued)





HIGH BYTE (DC8-DC15)

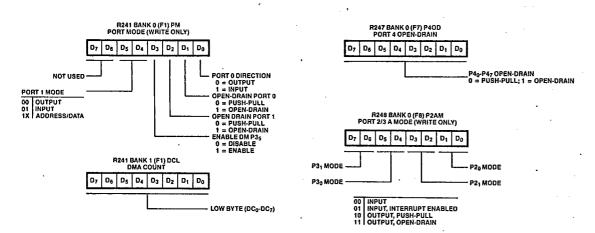
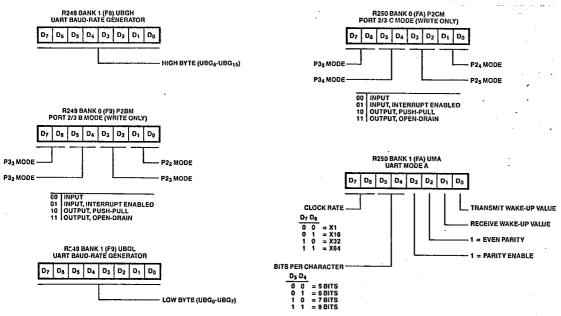
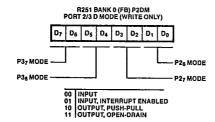


Figure 8. Mode and Control Registers (Continued)

MODE AND CONTROL REGISTERS (Continued)

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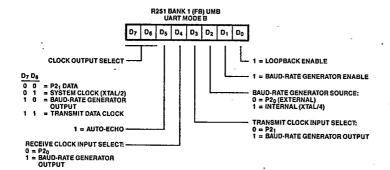
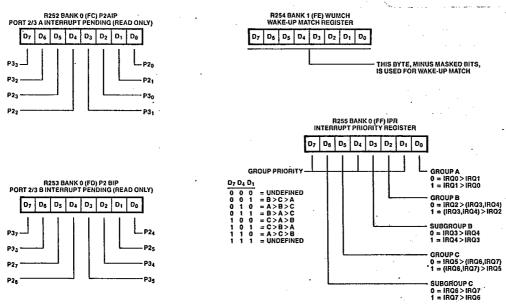
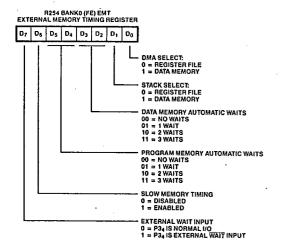


Figure 8. Mode and Control Registers (Continued)

MODE AND CONTROL REGISTERS (Continued)

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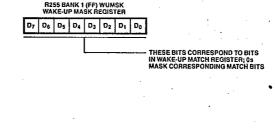


Figure 8. Mode and Control Registers (Continued)

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I/O PORTS

The Super8 has 40 I/O lines arranged into five 8-bit ports. These lines are all TTL-compatible, and can be configured as inputs or outputs. Some can also be configured as address/data lines.

Each port has an input register, an output register, and a register address. Data coming into the port is stored in the input register, and data to be written to a port is stored in the output register. Reading a port's register address returns the value in the input register; writing a port's register address loads the value in the output register. If the port is configured for an output, this value will appear on the external pins.

When the CPU reads the bits configured as outputs, the data on the external pins is returned. Under normal output loading, this has the same effect as reading the output register, unless the bits are configured as open-drain outputs.

The ports can be configured as shown in Table 2.

Table 2. Port Configuration

Port	Configuration Choices
0	Address outputs and/or general I/O
1	Multiplexed address/data(or I/O, only for ROM and Protopack)
2 and 3	Control I/O for UART, handshake channels, and counter/timers; also general I/O and external interrupts
4	General I/O

Port 0

Port 0 can be configured as an I/O port or an output for addressing external memory, or it can be divided and used as both. The bits configured as I/O can be either all outputs or all inputs; they cannot be mixed. If configured for outputs, they can be push-pull or open-drain type.

Any bits configured for I/O can be accessed via R208. To write to the port, specify R208 as the destination (dst) of an instruction; to read the port, specify R208 as the source (src).

Port 0 bits configured as I/O can be placed under handshake control of handshake channel 1.

Port 0 bits configured as address outputs cannot be accessed via the register.

In ROMless devices, initially the four lower bits are configured as address eight through twelve.

Port 1

In the ROMless device, Port 1 is configured as a byte-wide address/data port. It provides a byte-wide multiplexed address/data path. Additional address lines can be added by configuring Port 0.

The ROM and Protopack Port 1 can be configured as above or as an I/O port; it can be a byte-wide input, open-drain output, or push-pull output. It can be placed under handshake control or handshake channel 0.

Ports 2 and 3

Ports 2 and 3 provide external control inputs and outputs for the UART, handshake channels, and counter/timers. The pin assignments appear in Table 3.

Bits not used for control I/O can be configured as general-purpose I/O lines and/or external interrupt inputs.

Those bits configured for general I/O can be configured individually for input or output. Those configured for output can be individually configured for open-drain or push-pull output.

All Port 2 and 3 input pins are Schmitt-triggered.

The port address for Port 2 is R210, and for Port 3 is R211.

Table 3. Pin Assignments for Ports 2 and 3

Port	2	Port	3
Bit 0 1 2 3	Function	Bit	Function
0	UART receive clock	0	UART receive data
1	UART transmit clock	1	UART transmit data
2	Reserved	2	Reserved
3	Reserved	3	Reserved
4	Handshake 0 input	4	Handshake 1 input/WAIT
5	Handshake 0 output	5	Handshake 1 output/DM
6	Counter 0 input	6	Counter 1 input
7	Counter 0 I/O	7.	Counter 1 I/O

Port 4

Port 4 can be configured as I/O only. Each bit can be configured individually as input or output, with either push-pull or open-drain outputs. All Port 4 inputs are Schmitt-triggered.

Port 4 can be placed under handshake control of handshake channel 0. Its register address is R212.

UART

The UART is a full-duplex asynchronous channel. It transmits and receives independently with 5 to 8 bits per character, has options for even or odd bit parity, and a wake-up feature.

Data can be read into or out of the UART via R239, Bank 0. This single address is able to serve a full-duplex channel because it contains two complete 8-bit registers—one for the transmitter and the other for the receiver.

Pins

The UART uses the following Port 2 and 3 pins:

Port/Pin	UART Function
2/0	Receive Clock
3/0	Receive Data
2/1	Transmit Clock
3/1	Transmit Data

Transmitter

When the UART's register address is specified as the destination (dst) of an operation, the data is output on the UART, which automatically adds the start bit, the programmed parity bit, and the programmed number of stop bits. It can also add a wake-up bit if that option is selected.

If the UART is programmed for a 5-, 6-, or 7-bit character, the extra bits in R239 are ignored.

Serial data is transmitted at a rate equal to 1, 1/16, 1/32 or 1/64 of the transmitter clock rate, depending on the programmed data rate. All data is sent out on the falling edge of the clock input.

When the UART has no data to send, it holds the output marking (High). It may be programmed with the Send Break command to hold the output Low (Spacing), which it continues until the command is cleared.

Receiver

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The UART begins receive operation when Receive Enable (URC, bit 0) is set High. After this, a Low on the receive input pin for longer than half a bit time is interpreted as a start bit. The UART samples the data on the input pin in the middle of each clock cycle until a complete byte is assembled. This is placed in the Receive Data register.

If the 1X clock mode is selected, external bit synchronization must be provided, and the input data is sampled on the rising edge of the clock.

For character lengths of less than eight bits, the UART inserts ones into the unused bits, and, if parity is enabled, the parity bit is not stripped. The data bits, extra ones, and the parity bit are placed in the UART Data register (UIO).

While the UART is assembling a byte in its input shift register, the CPU has time to service an interrupt and manipulate the data character in UIO.

Once a complete character is assembled, the UART checks it and performs the following:

- If it is an ASCII control character, the UART sets the Control Character status bit.
- It checks the wake-up settings and completes any indicated action.
- If parity is enabled, the UART checks to see if the calculated parity matches the programmed parity bit. If they do not match, it sets the Parity Error bit in URC (R236 Bank 0), which remains set until reset by software.
- It sets the Framing Error bit (URC, bit 4) if the character is assembled without any stop bits. This bit remains set until cleared by software.

Overrun errors occur when characters are received faster than they are read. That is, when the UART has assembled a complete character before the CPU has read the current character, the UART sets the Overrun Error bit (URC, bit 3), and the character currently in the receive buffer is lost.

The overrun bit remains set until cleared by software.

ADDRESS SPACE

The Super8 can access 64K bytes of program memory and 64K bytes of data memory. These spaces can be either combined or separate. If separate, they are controlled by the $\overline{\rm DM}$ line (Port P3₅), which selects data memory when Low and program memory when High.

Figure 9 shows the system memory space.

CPU Program Memory

Program memory occupies addresses 0 to 64K. External program memory, if present, is accessed by configuring Ports 0 and 1 as a memory interface.

The address/data lines are controlled by \overline{AS} , \overline{DS} and R/\overline{W} .

The first 32 program memory bytes are reserved for interrupt vectors; the lowest address available for user programs is 32 (decimal). This value is automatically loaded into the program counter after a hardware reset.

ROMless

Port 0 can be configured to provide from 0 to 8 additional address lines. Port 1 is always used as an 8-bit multiplexed address/data port.

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ROM and Protopack

Port 1 is configured as multiplexed address/data or as I/O. When Port 1 is configured as address/data, Port 0 lines can be used as additional address lines, up to address 15. External program memory is mapped above internal program memory; that is, external program memory can occupy any space beginning at the top of the internal ROM space up to the 64K (16-bit address) limit.

CPU Data Memory

The external CPU data memory space, if separated from program memory by the \overline{DM} optional output, can be mapped anywhere from 0 to 64K (full 16-bit address space). Data memory uses the same address/data bus (Port 1) and additional addresses (chosen from Port 0) as program memory. Data memory is distinguished from program memory by the DM pin (P3₅), and by the fact that data memory can begin at address 0000_H. This feature differs from the Z8.

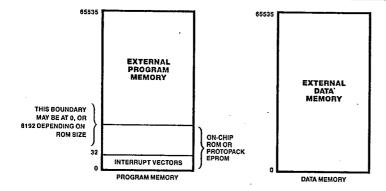


Figure 9. Program and Data Memory Address Spaces

INSTRUCTION SET

The Super8 instruction set is designed to handle its large register set. The instruction set provides a full complement of 8-bit arithmetic and logical operations, including multiply and divide. It supports BCD operations using a decimal adjustment of binary values, and it supports incrementing and decrementing 16-bit quantities for addresses and counters.

It provides extensive bit manipulation, and rotate and shift operations, and it requires no special I/O instructions—the I/O ports are mapped into the register file.

Instruction Pointer

A special register called the Instruction Pointer (IP) provides hardware support for threaded-code languages. It consists of register-pair R218 and R219, and it contains memory addresses. The MSB is R218.

Threaded-code languages deal with an imaginary higher-level machine within the existing hardware machine. The IP acts like the PC for that machine. The command NEXT passes control to or from the hardware machine to the imaginary machine, and the commands ENTER and EXIT are imaginary machine equivalents of (real machine) CALLS and RETURNS.

If the commands NEXT, ENTER, and EXIT are not used, the IP can be used by the fast interrupt processing, as described in the Interrupts section.

Flag Register

The Flag register (FLAGS) contains eight bits that describe the current status of the Super8. Four of these can be tested and used with conditional jump instructions; two others are used for BCD arithmetic. FLAGS also contains the Bank Address bit and the Fast Interrupt Status bit.

The flag bits can be set and reset by instructions.

CAUTION

Do not specify FLAGS as the destination of an instruction that normally affects the flag bits or the result will be unspecified.

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The following paragraphs describe each flag bit:

Bank Address. This bit is used to select one of the register banks (0 or 1) between (decimal) addresses 224 and 255. It is cleared by the SB0 instruction and set by the SB1 instruction.

Fast Interrupt Status. This bit is set during a fast interrupt cycle and reset during the IRET following interrupt servicing. When set, this bit inhibits all interrupts and causes the fast interrupt return to be executed when the IRET instruction is fetched.

Half-Carry. This bit is set to 1 whenever an addition generates a carry out of bit 3, or when a subtraction borrows out of bit 4. This bit is used by the Decimal Adjust (DA) instruction to convert the binary result of a previous addition or subtraction into the correct decimal (BCD) result. This flag, and the Decimal Adjust flag, are not usually accessed by users.

Decimal Adjust. This bit is used to specify what type of instruction was executed last during BCD operations, so a subsequent Decimal Adjust operation can function correctly. This bit is not usually accessible to programmers, and cannot be used as a test condition.

Overflow Flag. This flag is set to 1 when the result of a twos-complement operation was greater than 127 or less than -128. It is also cleared to 0 during logical operations.

Sign Flag. Following arithmetic, logical, rotate, or shift operations, this bit identifies the state of the MSB of the result. A 0 indicates a positive number and a 1 indicates a negative number.

Zero Flag. For arithmetic and logical operations, this flag is set to 1 if the result of the operation is zero.

For operations that test bits in a register, the zero bit is set to 1 if the result is zero.

For rotate and shift operations, this bit is set to 1 if the result is zero.

Carry Flag. This flag is set to 1 if the result from an arithmetic operation generates a carry out of, or a borrow into, bit 7.

After rotate and shift operations, it contains the last value shifted out of the specified register.

It can be set, cleared, or complemented by instructions.

Condition Codes

The flags C, Z, S, and V are used to control the operation of conditional jump instructions.

The opcode of a conditional jump contains a 4-bit field called the condition code (cc). This specifies under which conditions it is to execute the jump. For example, a conditional jump with the condition code for "equal" after a compare operation only jumps if the two operands are equal.

The condition codes and their meanings are given in Table 4.

Addressing Modes

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All operands except for immediate data and condition codes are expressed as register addresses, program memory addresses, or data memory addresses. The addressing modes and their designations are:

Register (R)
Indirect Register (IR)
Indexed (X)
Direct (DA)
Relative (RA)
Immediate (IM)
Indirect (IA)

Table 4. Condition Codes and Meanings

Binary	Mnemonic	Flags	Meaning .
0000	F	_	Always false
1000	_	-	Always true
0111*	С	C = 1	Carry
1111*	NC	C=0	No carry
0110*	Z	Z = 1	Zero
1110*	NZ	Z=0	Not zero
1101	PL	S=0	Plus
0101	MI	S = 1	Minus
0100	OV	V == 1	Overflow
1100	NOV	V = 0	No overflow
0110*	EQ	Z=1	Equal
1110*	NE	Z=0	Not equal
1001	GE	(S XOR V) = 0	Greater than or equal
0001	LT	(S XOR V) = 1	Less than
1010 •	GT	(Z OR (S XOR V)) = 0	Greater than
0010	LE	(Z OR (S XOR V)) = 1	Less than or equal
1111*	UGE	C=0	Unsigned greater than or equal
0111*	ULT	C=1	Unsigned less than
1011	UGT	(C = 0 AND Z = 0) ≈ 1	Unsigned less than
0011	ULE	(C OR Z) = 1	Unsigned greater than Unsigned less than or equal

NOTE: Asterisks (*) indicate condition codes that relate to two different mnemonics but test the same flags. For example, Z and EQ are both True if the Zero flag is set, but after an ADD instruction, Z would probably be used, while after a CP instruction, EQ would probably be used.

Registers can be addressed by an 8-bit address in the range of 0 to 255. Working registers can also be addressed using 4-bit addresses, where five bits contained in a register pointer (R218 or R219) are concatenated with three bits from the 4-bit address to form an 8-bit address.

Registers can be used in pairs to generate 16-bit program or data memory addresses.

Notation and Encoding

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The instruction set notations are described in Table 5.

Functional Summary of Commands

Figure 10 shows the formats followed by a quick reference guide to the commands.

Table 5. Instruction Set Notations

Notation	Meaning	Notation	Meaning
CC	Condition code (see Table 4)	DA	Direct address (between 0 and 65535)
ſ	Working register (between 0 and 15)	RA	Relative address
rb	Bit of working register	IM	Immediate
rO	Bit 0 of working register	IML	Immediate long
R	Register or working register	dst	Destination operand
RR	Register pair or working register pair (Register pairs	src	Source operand
•	always start on an even-number boundary)	@	Indirect address prefix
IA	Indirect address	SP	Stack pointer
lr.	Indirect working register	PC	Program counter
iR	Indirect register or indirect working register	IP	Instruction pointer
Irr	Indirect working register pair	FLAGS	Flags register
IRR	Indirect register pair or indirect working register pair	RP	Register pointer
X	Indexed	#	Immediate operand prefix
XS	Indexed, short offset	%	Hexadecimal number prefix
XL	Indexed, long offset	OPG	Opcode

One-Byte Instructions

	F, DI. EI. ENTER. EXIT. IRET, NEXT. NOP. F, RET. SBO, SB1. SCF, WFI
dst OPC INC	· ·
	Two-Byte Instructions
OPC	dst src ADC, ADD, AND, CP, LD, LDC, LDCI, LDCD, LDE, LDED, OR, SBC, SUB, TCM, TM, XOR
OPC	src dst LDC, LOCPD, LDCPI, LDE, LDEPD, LDEPI
OPC	CALL, DA, DEC, DECW. INC, INCW, JP, POP, RL, RLC, RR, RRC, SWAP, CLR, SRA, COM
OPC	STC PUSH, SRP, SRP0, SRP1
OPC	dst b 0 BITC, BITR
OPC	dsi b 1 BiTS
r OPC	dst DJNZ
cc OPC	dst JR
dst OPC	src LO
src OPC	dst LD

Figure 10. Instruction Formats

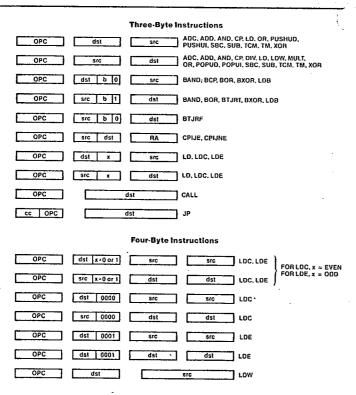


Figure 10. Instruction Formats (Continued)

INSTRUCTION SUMMARY

Instruction	Addr	Mode	Opcode	Flags Affected									
and Operation	dst	src	Byte (Hex)	С	Z	s	٧	D	Н				
ADC dst,src dst ← dst + src + C	(No	te 1)	1□	*	*	*	_	0	*				
ADD dst,src dst ← dst + src	(No	te 1)	0□	*	*	*	*	0	*				
AND dst,src . dst ← dst AND src	(No	te 1)	5□	_	*	*	0		_				
BAND dst,src dst ← dst AND src	r0 Rb	Rb r0	67 67	_	*	0	U	_	_				
BCP dst, src dst - src	r0	Rb	17		*	0	U	-	_				
BITC dst dst ← NOT dst	rb		57	-	*	0	U	_	_				
BITR dst dst ← 0	rb		77	_	_	_	_	_	_				
BITS dst dst ← 1	rb	 -	77	_	_	_	_	_	_				

Instruction	Addr	Mode	Opcode	Flags Affected								
and Operation	dst	src	Byte (Hex)	С	z	s	٧	D	Н			
BOR dst, src	r0	rB	07	_	*	0	U	_	_			
dst ← dst OR src	Rb	r0										
BTJRF	RA	rb	37			_	_		_			
if $src = 0$, $PC = PC$	+ dst					,						
BTJRT	RA	rb	37	_		÷	_	_				
if src = 1, PC = PC	+ dst											
BXOR dst, src	rO	Rb	27	_	*	0	U	_	_			
dst ← dst XOR src	Rb	r0	27									
CALL dst	DA		F6	_	_	_	_		_			
SP ← SP - 2	IRR		F4									
@SP ← PC	IA		D4									
PC ← dst												
CCF			EF	*	<u>-</u>	_		_	_			
C = NOT C												
CLR dst	R		B0	`	_	_	_	_				
dst ← 0	IR		B1									

INSTRUCTION SUMMARY (Continued)

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	Addr	Mode	Opcode	F	lag	s A	Aff	ec	tec	•		ddri	Vlode	•	Fla	gs	Af	fec	cted
Instruction and Operation	dst	src	Byte (Hex)	C	z	s	٧	, I	D I	H	Instruction and Operation	dst	src	Byte (Hex)	C	Z	S	٧	DН
COM dst dst ← NOT dst	R		60 61	_	*	*	0) -		_		RR IR		A0 A1	- 1	*	* :	k	
CP dst,src dst - src	(No	te 1)	Α□	*	*	*	*				IRET (Fast) PC ↔ IP			BF	Rest befo				ıpt
CPIJE if dst - src = 0,then	r	İr	C2	_	_	_				_	FLAG ← FLAG' FIS ← 0								· —
PC ← PC + RA Ir ← Ir + 1											IRET (Normal) FLAGS ← @SP; SP ← SP +			BF 0) ← 1	Resi				ıpt
CPIJNE	r	lr	D2	-	_	-	-			-		DA		ccD	<u> </u>				
if dst - src = 0,then PC ← PC + RA Ir ← Ir + 1											if cc is true,	IRR		(cc = 0 to F))				
DA dst	R		40	*	*	*		٠ ر	_	_	JR cc,dst	RA		ccB			_		
dst ← DA dst	IR R		41			_		_			if cc is true, PC ← PC + d			(cc = 0 to F)				
DEC dst dst ← dst – 1	IR		01		-	75		•			LD dst,src	r	IM	rC	_		_		
DECW dst	RR		80			_		<u> </u>		_	dst ← src	r	R	r8					
dst ← dst – 1	IR		81		_	_		•				R	r	r9 (r = 0 to F)	•				
DI			8F	-	-	-		-	_	-		r IR	IR r	C7 D7					
SMR (0) ← 0							_					R	Ŕ	. E4					
DIV dst, src		_									•	R	IR	E5					
dst ÷ src	RR	R	94	. *	*	*	*	*	_	_		R	IM	E6					
dst (Upper) ← Quotient	RR	IR	95									IR	IM	D6					
dst (Lower) ←	RR	IM	96									IR r	R x	F5 87					
Remainder										_		X	r	97					
DJNZ r,dst	RA	r	rA	-		-		_	_	-	LDB dst, src	r0	Rb	47			_		
r←r – 1			(r = 0 to F)								dst ← src	Rb	rO	47					
if r = 0 PC ← PC + dst											LDC/LDE	r	irr	СЗ		÷			
			9F				_	_	_	_	dst ← src	irr	r	D3					
EI SMR (0) ← 1			91	_	_				_			r	XS	E7					
												XS	г х1	F7 A7					
ENTER			1F	_				_				r ∙x1	r	B7					
SP ← SP - 2 @ SP ← IP	•											r	DA.			•			
IP ← PC -												DΑ	r	B7					
PC ← @ IP IP ← IP + 2											LDCD/LDED dst, src dst ← src	r	Irr	E2	_	_		-	
EXIT			2F	-		-	_	_	_	_	rr ← rr − 1								
IP ← @SP SP ← SP + 2											LDEI/LDCI dst, src dst ← src	r	İrr	E3	_	_		-	
PC ← @IP IP ← IP + 2											rr ← rr + 1								-
INC dst	r		rE			_	*	+	_	_	LDCPD/LDEPD dst,s								
dst ← dst + 1	'		(r = 0 to F)	•	• '	-	•			rr ← rr – 1 '	irr	r	F2	_	_		_	
	R		20	•							dst ← src				<u>.</u>		-		
	IR		21																

INSTRUCTION SUMMARY (Continued)

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Instruction	Addr	Mode	Opcode Byte	Flags Affected					
and Operation	dst	src	(Hex)	С	z	s	٧	D	Н
LDCPI/LDEPI dst, sre					_		_		
rr ← rr + 1 dst ← src	irr	r	F3	_	_	_		_	-
LDW dst, src	RR	RR	C4	_	_	_	_	_	_
dst ← src	RR	IR	C5						
	RR	IMM	C6						_
MULT dst, src	RR RR	R IR	84 85	*	0	*	*	_	-
	RR	IM	86						
NEXT			OF.		_	_	_	_	_
PC ← @IP			٠.					•	
IP ← IP + 2									
NOP			FF		_	_	_	_	_
OR dst,src dst ← dst OR src	(No	te 1)	4□	_	*	*	0	-	_
POP dst		R	50	_	_	_	_	_	_
dst ← @SP; SP ← SP + 1		IR	51	•					
POPUD dst, src dst ← src	R	IR	92	_	_		_	_	_
IR ← IR ~ 1									
POPUI dst, src	R	IR	93			_		_	_
dst ← src									
IR ← IR + 1									
PUSH src		R	70	_	_	_		_	-
SP ← SP - 1; @SP ←	src	IR	71						
PUSHUD dst, src IR ← IR – 1 dst ← src	IR	R	82		_	_	_	_	
PUSHUI dst, src IR ← IR + 1 dst ← src	IR	R	83	_	_	_	_	_	_
RCF C ← 0		- <u> </u>	CF	0		_	_	_	_
RET PC ← @SP; SP ← SP	+ 2		AF	_		_		_	_
RL dst	R		90	*	*	*	*	_	_
C ← dst (7) dst (0) ← dst (7) dst (N + 1) ← dst (N) N = 0 to 6	IR		91	٠					

					٠				
Instruction	Addr	Mode	Opcode Byte	F	lag	s A	Affe	ecte	₽d
and Operation	dst	src	(Hex)	C	Z	S	٧	Ď	н
RLC dst dst (0) \leftarrow C C \leftarrow dst (7) dst (N + 1) \leftarrow dst (N) N = 0 to 6	R IR		10 11	*	*	*	*		
RR dst $C \leftarrow dst(0)$ $dst(7) \leftarrow dst(0)$ $dst(N) \leftarrow dst(N + 1)$ N = 0 to 6	R		E0 E1	*	*	*	*		_
RRC dst C ← dst (0) dst (7) ← C dst (N) ← dst (N + 1) N = 0 to 6	R IR		C0 C1	*	*	*	*	-	.—
SB0 BANK ← 0			4F		_	_	_	_	_
SB1 BANK ← 1			5F	_	-	_	_	_	_
SBC dst,src dst ← dst – src – C	(No	te 1)	3□	*	*	*	*	1	*
SCF C ← 1			DF	1	_	-	_		_
SRA dst dst (7) ← dst (7) C ← dst (0) dst (N) ← dst (N + 1) N = 0 to 6	R IR		D0 D1	*	*	*	0	_	
SRP src RP0 ← IM RP1 ← IM + 8		IM	31	_	_	_		_	_
SRP0 RP0 ← IM		IM	31	_	_	_		_	
SRP1 RP1 ←IM		IM	- 31	_	_		-	_	_
SUB dst,src dst ← dst – src	(Not	e 1)	2□	*	*	*	*	1	*
					_		_		

INSTRUCTION SUMMARY (Continued)

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Table 6. Second Nibble

	Addr	Mode	Opcode	F	lag	s A	ffe	cte	d
Instruction and Operation	dst src		Byte (Hex)	C	z	s	٧	D	Н
SWAP dst dst (0·3) ++ dst (4·7)	R IR		F0 F1	_	*	*	U	_	_
TCM dst,src (NOT dst) AND src	(No	te 1)	6□	_	*	*	0	_	_
TM dst,src dst AND src	(No	te 1)	7□	_	*	*	0	_	_
WFI			3F	_	_	_	_	_	
XOR dst,src dst ← dst XOR src	(No	ote 1)	B□		*	*	0		_

NOTE 1: These instructions have an identical set of addressing modes, which are encoded for brevity. The first opcode nibble identifies the command, and is found in the table above. The second nibble, represented by a □, defines the addressing mode as shown in Table 6.:

Addr	Mode	Lower
dst	src	Opcode Nibble
r	r	2
r	lr	3
R	R	4
R	` IR	5
R	IM	6

For example, to use an opcode represented as x□ with an "RR" addressing mode, use the opcode "x4."

- 0 = Cleared to Zero
- 1 = Set to One
- = Unaffected
- U = Undefined

SUPER-8 OPCODE MAP

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								Low	er Nibb	le (He	x)								
		0	1	2	3	4	5	6	7	8		9	A	В	C	D		E	F
	0	DEC	6 DEC IR ₁	6 ADD (1./2	ADD r ₁ ,lr ₂	10 ADD R ₂ ,R ₁	10 ADD IR ₂ ,R ₁	10 ADD R ₁ ,IM	10 BOR*	6 LD r ₁ ,R ₂	- -	6 LD 2.R ₁	12/10 DJNZ r ₁ ,RA	12/10 JR cc,RA	6 LD (1,1M	12/1 JP cc,D		6 INC r1	14 NEXT
	1	6 RLC R ₁	6 RLC IR ₁	6 ADC r ₁ ,r ₂	6 ADC r ₁ ,lr ₂	10 ADC R ₂ ,R ₁	10 ADC IR ₂ R ₁	10 ADC R ₁ .IM	10 BCP r ₁ ,b.R ₂									Ï	20 ENTER
	2	6 INC R ₁	6 INC IR ₁	6 SUB r ₁ ,r ₂	6 SUB (1.l(2	10 SUB R ₂ ,R ₁	10 SUB IR ₂ ,R ₁	10 SUB R ₁ .IM	10 BXOR*										22 EXIT
	3	10 JP IRR ₁	NOTE C	6 SBC r ₁ .r ₂	6 SBC r ₁ ,lr ₂	10 SBC R ₂ ,R ₁	10 SBC IR ₂ ,R ₁	10 SBC R ₁ .IM	NOTE A										6 WFI
	4	6 DA R ₁	6 DA IR ₁	6 OR (1./2	6 OR r ₁ ,lr ₂	10 OR R ₂ ,R ₁	10 OR IR ₂ ,R ₁	10 OR R ₁ .iM	10 LDB*										6 SBO
	5	POP R ₁	10 POP IR ₁	6 AND 11.12	6 AND r ₁ .lr ₂	10 AND R ₂ .R ₁	10 AND IR ₂ ,R ₁	10 AND R ₁ .IM	8 BITC r ₁ ,b										6 SBI
Hex)	6	COM R ₁	COM	6 TCM 11.12	6 TCM (1.lr ₂	10 TCM R ₂ ,R ₁	10 TCM IR ₂ ,R ₁	10 TCM R ₁ .IM	10 BAND* r ₀ ·R _b										
libble (I	7	10:12 PUSH R ₂	12/14 PUSH IR ₂	6 TM (1./2	6 TM rt,ir2	10 TM R ₂ .R ₁	10 TM IR ₂ ,R ₁	10 TM R ₁ .IM	NOTE B										
Upper Nibble (Hex)	8	10 DECW RR ₁	10 DECW IR ₁	10 PUSHUD IR ₁ .R ₂	10 PUSHUI IR ₁ .R ₂	24 MULT R ₂ ,RR ₁	24 MULT IR ₂ ,RR ₁	24 MULT IM.RR;	10 LD r ₁ .x.r ₂			į							6 DI
_	9	6 RL R ₁	6 RL IR ₁	10 POPUD IR _{2.} R ₁	10 POPUI IR ₂ ,R ₁	28/12 DIV R ₂ ,RR ₁	28/12 DIV IR ₂ ,RR ₁	28/12 DIV IM,RR ₁	10 LD r ₂ .x,r ₁		1								6 El
	A	10 INCW RR ₁	10 INCW IR ₁	6 CP r _{1.} r ₂	6 CP r ₁ .lr ₂	10 CP R ₂ .R ₁	10 CP IR ₂ ,R ₁	10 CP R ₁ ,IM	NOTE D										14 RET
	В	6 CLR R ₁	6 CLR IR,	6 XOR (1./2	6 XOR r ₁ .ir ₂	10 XOR R ₂ .R ₁	10 XOR IR ₂ ,R ₁	10 XOR R ₁ ,IM	NOTE E										16/6 IRET
	С	6 RAC R ₁	6 RRC IR ₁	16/18 CPIJE Icr ₂ ,RA	11'ju5 FDC.	10 LDW RR ₂ ,RR ₁	10 LDW IR ₂ .RR ₁	12 LOW RR ₁ .IML	6 LD r _{1.} lr ₂										6 RCF
	D	6 SRA R ₁	6 SRA IR ₁	16/18 CPIJNE Ir ₁ .r ₂ .RA	12.lt.1	20 CALL IA ₁		10 LD IR ₁ ,IM	/ 6 LD lr ₁ ,r ₂							ŀ			6 SCF
	E	6 RA R ₁	6 RR IR ₁	11'ju5 rdcd.	16 LDCI* r _{1.} lrr ₂	10 LD R ₂ ,R ₁	10 LD IR ₂ ,R ₁	10 LD R ₁ ,IM	18 LDC* r ₁ ,lrr ₂ ,xs										6 CCF
	F	8 SWAP R ₁	8 SWAP IR ₁	16 LDCPD*	16 LDCPI* r ₂ .lrr ₁	18 CALL IRR ₁	10 LD R ₂ ,IR ₁	18 CALL DA ₁	18 LDC* r ₂ .lrr ₁ .xs	<u> </u>							,		6 NOP

16/18 16/18 BTJRF BTJRT r₂.b RA r₂.b.RA NOTE A

8 BITR 8 BITS r₁.b NOTE B

NOTE C

6 SRP IM 6 SRP0 IM 6 SRP1 IM

NOTE D

20 LDC* r₁.lrr₂.xL 20 LDC* r₁.DA₂

NOTEE

20 LDC* r₂.DA₁ 20 LDC* r2,lrr2,xL

Legend:
r = 4-bit address
R = 8-bit address
b = bit number
R₁ or r₁ = dst address
R₂ or r₂ = src address

*Examples:
BOR r₀-R₂
is BOR r₁, b.R₂
or BOR r₂-b,R₁
LDC r₁.Irr₂ = program
or LDE r₁.Irr₂ = data

Sequence: Opcode, first, second, third operands

NOTE: The blank areas are not defined.

Figure 11. Opcode Map

INSTRUCTIONS

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Table 7. Super8 Instructions

Vinemonic	Operands	Instruction	Mnemonic	Operands	Instruction
Load Instruc	ctions		Program Co	ntrol Instructio	
CLR	dst	Clear	BTJRT	dst, src	Bit test jump relative on True
LD	dst, src	Load	BTJRF	dst, src	Bit test jump relative on False
_DB	dst, src	Load bit	CALL	dst	Call procedure
_DC	dst, src	Load program memory	CPIJE	dst, src	Compare, increment and jump on
LDE	dst, src	Load data memory	_		equal
LDCD	dst, src	Load program memory and	CPIJNE	dst, src	Compare, increment and jump on
_000	031, 310	decrement			non-equal
LDED	dst, src	Load data memory and	DJNŽ	r, dst	Decrement and jump on non-zero
LDLO	dat, are	decrement	ENTER		Enter
_DCI	dst, src	Load program memory and	EXIT		Exit
_DC1	dat, are	increment	IRET	·	Return from interrupt
חבו	dat ara	Load data memory and increment	JP	cc, dst	Jump on condition code
LDEI	dst, src	Load program memory with	JP	dst	Jump unconditional
LDCPD	dst, src	pre-decrement	JR	cc, dst	Jump relative on condition code
. הבהה	dat ara	· ·	JR	dst	Jump relative unconditional
LDEPD	dst, src	Load data memory with	NEXT		Next
		pre-decrement	RET		Return
LDCPI	dst, src	Load program memory with pre-increment	WFI		Wait for interrupt
LDEPI	dst, src	Load data memory with	Bit Manipul	ation Instructi	ons
		pre-increment	BAND	dst, src	Bit AND
LDW	dst, src	Load word	BCP	dst, src	Bit compare
POP	dst	Pop stack	BITC	dst	Bit complement
POPUD	dst, src	Pop user stack (decrement)	BITR	dst	Bit reset
POPUI	dst, src	Pop user stack (increment)	BITS	dst	Bit set
PUSH	SIC	Push stack	BOR	dst, src	Bit OR
PUSHUD	dst, src	Push user stack (decrement)	BXOR	. det, ere	Bit exclusive OR
PUSHUI	dst, src	Push user stack (increment)	TCM	dst, src	Test complement under mask
			TM	dst, src	Test under mask
Arithmetic	Instructions		Rotate and	Shift Instructi	ions
ADC	dst, src	Add with carry	RL	dst	Rotate left
ADD	dst, src	Add	RLC	dst	Rotate left through carry
CP	dst, src	Compare	RR	dst	Rotate right
DA	dst	Decimal adjust	RRC	dst	Rotate right through carry
DEC	dst	Decrement	SRA	dst	Shift right arithmetic
DECW	dst	Decrement word .	SWAP	dst	Swap nibbles
DIV	dst, src	Divide			<u> </u>
INC	dst	Increment		ol Instructions	
INCW	dst	Increment word	CCF		Complement carry flag
MULT	dst, src	Multiply	DI		Disable interrupts
SBC	dst, src	Subtract with carry	El		Enable interrupts
SUB	dst, src	Subtract	NOP		Do nothing
			RCF		Reset carry flag
1 1 - 11			SB0		Set bank 0
Logical Ins	structions		SB1		Set bank 1 ,
AND	dst, src	Logical AND	SCF	٠	Set carry flag
	dst	Complement	SRP	src	Set register pointers
СОМ	Qat	· •- ·			
	dst, src	Logical OR	SRPO	src	Set register pointer zero

1002

INTERRUPTS

The Super8 interrupt structure contains 8 levels of interrupt, 16 vectors, and 27 sources.

Interrupt priority is assigned by level, controlled by the Interrupt Priority register (IPR). Each level is masked (or enabled) according to the bits in the Interrupt Mask register (IMR), and the entire interrupt structure can be disabled by clearing a bit in the System Mode register (R222).

The three major components of the interrupt structure are sources, vectors, and levels. These are shown in Figure 10 and discussed in the following paragraphs.

Sources

A source is anything that generates an interrupt. This can be internal or external to the Super8 MCU. Internal sources are hardwired to a particular vector and level, while external sources can be assigned to various external events.

External interrupts are falling-edge triggered.

Vectors

The 16 vectors are divided unequally among the eight levels. For example, vector 12 belongs to level 2, while level 3 contains vectors 0, 2, 4, and 6.

The vector number is used to generate the address of a particular interrupt servicing routine; therefore all interrupts using the same vector must use the same interrupt handling routine.

Levels

Levels provide the top level of priority assignment. While the sources and vectors are hardwired within each level, the priorities of the levels can be changed by using the Interrupt Priority register (see Figure 8 for bit details).

If more than one interrupt source is active, the source from the highest priority level will be serviced first. If both sources are from the same level, the source with the lowest vector will have priority. For example, if the UART Receive Data bit and UART Parity Error bit are both active, the UART Parity Error bit will be serviced first because it is vector 16, and UART receive data is vector 20.

The levels are shown in Figure 12. . .

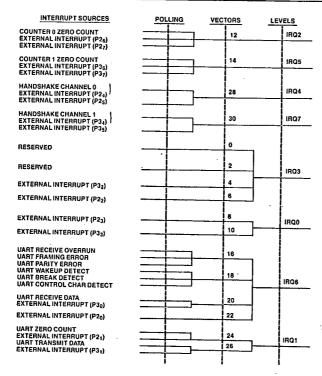


Figure 12. Interrupt Levels and Vectors

Enables

Interrupts can be enabled or disabled as follows:

- Interrupt enable/disable. The entire interrupt structure can be enabled or disabled by setting bit 0 in the System Mode register (R222).
- Level enable. Each level can be enabled or disabled by setting the appropriate bit in the Interrupt Mask register (R221).
- Level priority. The priority of each level can be controlled by the values in the Interrupt Priority register (R255, Bank 0).
- Source enable/disable. Each interrupt source can be enabled or disabled in the sources' Mode and Control register.

Service Routines

Before an interrupt request can be granted, a) interrupts must be enabled, b) the level must be enabled, c) it must be the highest priority interrupting level, d) it must be enabled at the interrupting source, and e) it must have the highest priority within the level.

If all this occurs, an interrupt request is granted.

The Super8 then enters an interrupt machine cycle that completes the following sequence:

- It resets the Interrupt Enable bit to disable all subsequent interrupts.
- It saves the Program Counter and status flags on the stack.
- It branches to the address contained within the vector location for the interrupt.
- It passes control to the interrupt servicing routine.

When the interrupt servicing routine has serviced the interrupt, it should issue an interrupt return (IRET) instruction. This restores the Program Counter and status flags and sets the Interrupt Enable bit in the System Mode register.

Fast Interrupt Processing

The Super8 provides a feature called fast interrupt processing, which completes the interrupt servicing in 6 clock periods instead of the usual 22.

Two hardware registers support fast interrupts. The Instruction Pointer (IP) holds the starting address of the service routine, and saves the PC value when a fast interrupt occurs. A dedicated register, FLAG', saves the contents of the FLAGS register when a fast interrupt occurs.

To use this feature, load the address of the service routine in the Instruction Pointer, load the level number into the Fast Interrupt Select field, and turn on the Fast Interrupt Enable bit in the System Mode register.

When an interrupt occurs in the level selected for fast interrupt processing, the following occurs:

- The contents of the Instruction Pointer and Program Counter are swapped.
- The contents of the Flag register are copied into FLAG!
- The Fast Interrupt Status Bit in FLAGS is set.
- The interrupt is serviced.

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- When IRET is issued after the interrupt service outline is completed, the Instruction Pointer and Program Counter are swapped again.
- The contents of FLAG' are copied back into the Flag register.
- The Fast Interrupt Status bit in FLAGS is cleared.

The interrupt servicing routine selected for fast processing should be written so that the location after the IRET instruction is the entry point the next time the (same) routine is used.

Level or Edge Triggered

Because internal interrupt requests are levels and interrupt requests from the outside are (usually) edges, the hardware for external interrupts uses edge-triggered flip-flops to convert the edges to levels.

The level-activated system requires that interrupt-serving software perform some action to remove the interrupting source. The action involved in serving the interrupt may remove the source, or the software may have to actually reset the flip-flops by writing to the corresponding Interrupt Pending register.

STACK OPERATION

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The Super8 architecture supports stack operations in the register file or in data memory. Bit 1 in the external Memory Timing register (R254 bank 0) selects between the two.

Register pair 216-217 forms the Stack Pointer used for all stack operations. R216 is the MSB and R217 is the LSB.

The Stack Pointer always points to data stored on the top of the stack. The address is decremented prior to a PUSH and incremented after a POP.

The stack is also used as a return stack for CALLs and interrupts. During a CALL, the contents of the PC are saved on the stack, to be restored later. Interrupts cause the contents of the PC and FLAGS to be saved on the stack, for recovery by IRET when the interrupt is finished.

When the Super8 is configured for an internal stack (using the register file), R217 contains the Stack Pointer. R216 may be used as a general purpose register, but its contents will be changed if an overflow or underflow occurs as the result of incrementing or decrementing the stack address during normal stack operations.

User-Defined Stacks

The Super8 provides for user-defined stacks in both the register file and program or data memory. These can be made to increment or decrement on a push by the choice of opcodes. For example, to implement a stack that grows from low addresses to high addresses in the register file, use PUSHUI and POPUD. For a stack that grows from high addresses to low addresses in data memory, use LDEI for pop and LDEPD for push.

COUNTER/TIMERS

The Super8 has two identical independently programmable 16-bit counter/timers that can be cascaded to produce a single 32-bit counter. They can be used to count external events, or they can obtain their input internally. The internal input is obtained by dividing the crystal frequency by four.

The counter/timers can be set to count up or down, by software or external events. They can be set for single or continuous cycle counting, and they can be set with a bi-value option, where two preset time constants alternate in loading the counter each time it reaches zero. This can be used to produce an output pulse train with a variable duty cycle.

The counter/timers can also be programmed to capture the count value at an external event or generate an interrupt whenever the count reaches zero. They can be turned on and off in response to external events by using a gate and/or a trigger option. The gate option enables counts only when the gate line is Low; the trigger option turns on the counter after a transient High. The gate and trigger options used together cause the counter/timer to work in gate mode after initially being triggered.

The control and status register bits for the counter/timers are shown in Figure 5.

DMA

The Super8 features an on-chip Direct Memory Access (DMA) channel to provide high bandwidth data transmission capabilities. The DMA channel can be used by the UART receiver, UART transmitter, or handshake channel 0. Data can be transferred between the peripheral and contiguous locations in either the register file or external

data memory. A 16-bit count register determines the number of transactions to be performed; an interrupt can be generated when the count is exhausted. DMA transfers to or from the register file require six CPU clock cycles; DMA transfers to or from external memory take ten CPU clock cycles, excluding wait states.

ABSOLUTE MAXIMUM RATINGS

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Voltage on all pins with respect
to ground 0.3V to +7.0V
Ambient Operating
Temperature See Ordering Information
Storage Temperature 65 °C to + 150 °C

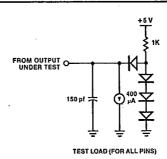
Stresses greater than these may cause permanent damage to the device. This is a stress rating only; operation of the device under conditions more severe than those listed for operating conditions may cause permanent damage to the device. Exposure to absolute maximum ratings for extended periods may also cause permanent damage.

STANDARD TEST CONDITIONS

Figure 14 shows the setup for standard test conditions. All voltages are referenced to ground, and positive current flows into the reference pin.

Standard conditions are:

- $+4.75V \le V_{CC} \le +5.25V$
- GND = 0V
- \blacksquare 0°C \leq T_A \leq +70°C



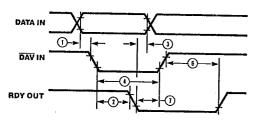
Standard Test Load

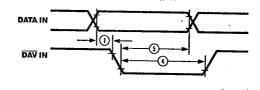
DC CHARACTERISTICS

Symbol	Parameter	Min	Max	Unit	Condition
V _{CH}	Clock Input High Voltage	3.8	Vcc	٧	Driven by External Clock Generator
VCL	Clock Input Low Voltage	-0.3	8.0	٧	Driven by External Clock Generator
V _{IH}	Input High Voltage	2.2	Vcc	٧	
V _{IL}	Input Low Voltage	-0.3	0.8	٧	
V _{RH}	Reset Input High Voltage	3.8	Vcc	٧	
V _{RL}	Reset Input Low Voltage	-0.3	8.0	٧	
VoH	Output High Voltage	2.4		٧	$I_{OH} = -400 \mu\text{A}$
VOL	Output Low Voltage		0.4	٧	$l_{OL} = +4.0 \text{mA}$
4L	Input Leakage	- 10	10	μΑ	
loL	Output Leakage	-10	10	μΑ	
IR	Reset Input Current		-50	μΑ	
lcc	V _{CC} Supply Current		320	mA	

INPUT HANDSHAKE TIMING

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Fully Interlocked Mode

Strobed Mode

AC CHARACTERISTICS (20 MHz)

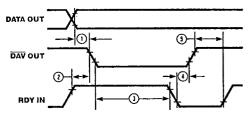
Input Handshake

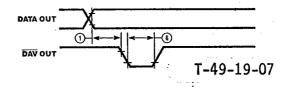
Number	Symbol	Parameter	Min	Max	Notes*#
1	TsDf(DAV)	Data In to Setup Time	0		
2	TdDAVif(RDY)	DAV Input to RDY I Delay		200	4
3	ThDI(RDY)	Data In Hold Time from RDY ↓	0	,	
4	TwDAV	DAV In Width	45		
5	ThDI(DAV)	Data In Hold Time from DAV ↓	130		
6	TdDAV(RDY)	DAV † Input to RDY † Delay	,50	100	. 0
7	TdRDYf(DAV)	RDY ↓ Output to DAV † Delay	0	. 100	2

NOTES:

Standard Test Load
 This time assumes user program reads data before DAV Input goes high. RDY will not go high before data is read. ‡Times given are in ns.
•Times are preliminary and subject to change.

OUTPUT HANDSHAKE TIMING





Fully Interlocked Mode

Strobed Mode

AC CHARACTERISTICS (12 MHz, 20 MHz)

Output Handshake

Number	Symbol	Parameter	Min	Max	Notes*‡
1	TdDO(DAV)	Data Out to DAV ↓ Delay	90		1,2
2	TdRDYr(DAV)	RDY † Input to DAV ↓ Delay	0	. 110	. 1
3	TdDAVOf(RDY)	DAV ↓ Output to RDY ↓ Delay	0		
4	TdRDYf(DAV)	RDY Input to DAV ↑ Delay	0	110	1
5	TdDAVOr(RDY)	DAV † Output to RDY † Delay	0 .		
6	TwDAVO	DAV Output Width	150		2

NOTES

- 2. Time given is for zero value in Deskew Counter. For nonzero value of n where n = 1, 2, ... 15 add $2 \times n \times TpC$ to the given time.
- ‡Times given are in ns.
 *Times are preliminary and subject to change.

AC CHARACTERISTICS (12 MHz)

Read/Write

			Norma	lTiming	Extende	d Timing	
Number	Symbol	Parameter	Min	Max	Min	Max	Notes‡*
1	TdA(AS)	Address Valid to AS † Delay	35		115		
2	TdAS(A)	AS t to Address Float Delay	65		150 .		
3	TdAS(DR)	AS t to Read Data Required Valid		270		600	1
4	TwAS	AS Low Width	65		150		
5	TdA(DS)	Address Float to DS ↓	20		20		
6a	TwDS(Read)	DS (Read) Low Width	225		470		1
6b	TwDS(Write)	DS (Write) Low Width	130		295		1
7	TdDS(DR)	DS I to Read Data Required Valid		180		420	.1
8	ThDS(DR)	Read Data to DS t Hold Time	0		0	-	
9	TdDS(A)	DS t to Address Active Delay	50	•	135	•	
10	TdDS(AS)	DS t to AS ↓ Delay	60		145		:
11	TdDO(DS)	Write Data Valid to DS (Write) ↓ Delay	35		115		
12	TdAS(W)	ĀS ↑ to Wait Delay		220		600	2
13	ThDS(W)	DS t to Wait Hold Time	0		0		
14	TdRW(AS)	R/W Valid to AS t Delay	50		135		

NOTES:

- 1. WAIT states add 167 ns to these times.
- 2. Auto-wait states add 167 ns to this time.
- All times are in ns and are for 12 MHz input frequency.
 Timings are preliminary and subject to change.

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AC CHARACTERISTICS (20 MHz)

Read/Write

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		·	Noma	Timing	Extende	d Timing		
Number	Symbol	Parameter	Min	Max	Min	Max	Notes‡	
1	TdA(AS)	Address Valid to AS † Delay	10		50			
2	TdAS(A)	AS t to Address Float Delay	35		85			
3	TdAS(DR)	AS to Read Data Required Valid		140		335	1	
4	TwAS	AS Low Width	35		85			
5	TdA(DS)	Address Float to DS ↓	0		0			
6a	TwDS(Read)	DS (Read) Low Width	125		275		1	
6b	TwDS(Write)	DS (Write) Low Width	65		165		1	
7	TdDS(DR)	DS ↓ to Read Data Required Valid		80		225	1	
8	ThDS(DR)	Read Data to DS t Hold Time	0		0			
9	TdDS(A)	DS t to Address Active Delay	20		70			
10	TdDS(AS)	DS † to AS ↓ Delay	30	1.	80			
11	TdDO(DS)	Write Data Valid to DS (Write) ↓ Delay	10		50			
12	TdAS(W)	ĀS † to Wait Delay	-	90		335	2	
13	ThDS(W)	DS t to Wait Hold Time	0		0		_	
14	TdRW(AS)	R/₩ Valid to ĀS † Delay	20	-	70			
15	TdDS(DW)	, DS t to Write Data Not Valid Delay	20		70			

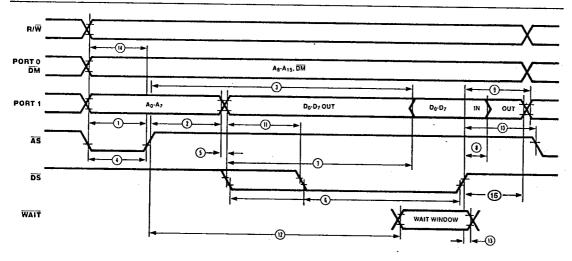
- NOTES:

 1. WAIT states add 100 ns to these times.

 2. Auto-wait states add 100 ns to this time.

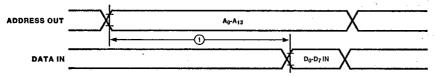
 ‡ All times are in ns and are for 20 MHz input frequency.

 Timings are preliminary and subject to change.



External Memory Read and Write Timing

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EPROM Read Timing

AC CHARACTERISTICS (20 MHz) EPROM Read Cycle

Number	Symbol	Parameter	Min	Max	Notes‡*
1	TdA(DR)	Address Valid to Read Data Required			,
		Valid		170	* 1 -

NOTES:

1. WAIT states add 167 ns to these times.

‡All times are in ns and are for 12 MHz input frequency.

*Timings are preliminary and subject to change.